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Validated polynomial optimisation for computer-assisted proofs in ODEs

Jeremy Parker

Motivating example: Lyapunov functions

Consider an autonomous ODE system

$$\frac{d}{dt}x(t) = f(x(t)), \quad f \in C^1(\Omega, \mathbb{R}^n),$$

where

$$\Omega \subset \mathbb{R}^n$$

is forward invariant.

Suppose that $f(0) = 0$. If there exists $V \in C^1(\Omega, \mathbb{R})$ such that

$$\begin{aligned} \mathcal{L}_f V(x) &\leq 0, & \forall x \in \Omega, \\ V(x) &\geq |x|^2, & \forall x \in \Omega. \end{aligned}$$

Then the equilibrium $x = 0$ is globally (Lyapunov) stable in Ω .

Here $\mathcal{L}_f V = f \cdot \nabla V$ is the Lie derivative of V with respect to the flow.



Motivating example: Lyapunov functions

Consider an autonomous ODE system

$$\frac{d}{dt}x(t) = f(x(t)), \quad x \in \mathbb{R}^n, \quad f \in \mathbb{R}[x]^n,$$

where

$$\Omega = \{x \in \mathbb{R}^n \mid g(x) \geq 0\},$$

$$g \in \mathbb{R}[x],$$

Multivariate
polynomial in x with
real coefficients

is forward invariant.

Suppose that $f(0) = 0$. If there exists $V \in \mathbb{R}[x]$ such that

$$\begin{aligned} -\mathcal{L}_f V &= \sigma_0 + g \sigma_1, & \sigma_0, \sigma_1 &\in SOS, \\ V(x) - |x|^2 &= \sigma_2 + g \sigma_3, & \sigma_2, \sigma_3 &\in SOS, \end{aligned}$$

Then the equilibrium $x = 0$ is globally (Lyapunov) stable in Ω .

Here $\sigma \in SOS$ means σ is a polynomial sum-of-squares, i.e. $\sigma = \sum p_i^2$ for some $p_i \in \mathbb{R}[x]$.

SOS as SDP

$\sigma = \sum p_i^2$ for some $p_i \in \mathbb{R}[x]$ if and only if

$$\sigma = \mathbf{b}^T Q \mathbf{b}, \quad Q \succcurlyeq 0,$$

Q is a real symmetric
positive semidefinite
matrix

where $\mathbf{b}$ is a vector of monomials of sufficiently high degree.

So the problem reduces to finding $V \in \mathbb{R}[x]$ and $Q_1, Q_2, Q_3, Q_4 \succcurlyeq 0$ such that

$$\begin{aligned} -\mathcal{L}_f V &= \mathbf{b}^T Q_1 \mathbf{b} + g \mathbf{b}^T Q_2 \mathbf{b}, \\ V(x) - |x|^2 &= \mathbf{b}^T Q_3 \mathbf{b} + g \mathbf{b}^T Q_4 \mathbf{b}. \end{aligned}$$

This is a semidefinite program (SDP). Lots of floating point solvers exist.

Other Lyapunov/auxiliary/storage function methods

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Bounds for Deterministic and Stochastic Dynamical Systems using Sum-of-Squares Optimization*

G. Fantuzzi[†], D. Goluskin[‡], D. Huang[§], and S. I. Chernyshenko[†]

$$\mathbf{f} \cdot \nabla V + \varphi - L \geq 0 \quad \forall \mathbf{x} \in \mathbb{R}^n.$$

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Bounding Extreme Events in Nonlinear Dynamics Using Convex Optimization*

Giovanni Fantuzzi[†] and David Goluskin[‡]

$$\begin{aligned} \mathcal{L}V(t, x) &\leq 0 \quad \forall (t, x) \in \Omega, \\ \Phi(t, x) - V(t, x) &\leq 0 \quad \forall (t, x) \in \Omega. \end{aligned}$$

IOP Publishing | London Mathematical Society

Nonlinearity

Nonlinearity **33** (2020) 4878–4899

<https://doi.org/10.1088/1361-6544/ab8f7b>

Bounding extrema over global attractors using polynomial optimisation

David Goluskin[©]

$$\Phi_A^+ \leq \inf_{\substack{\lambda > \mathbb{R} \\ V \in \mathcal{C}^1}} C \quad \text{s.t.} \quad \begin{aligned} V(\mathbf{x}) - \Phi(\mathbf{x}) &\geq 0 \quad \forall \mathbf{x} \in \mathbb{R}^n, \\ C - V(\mathbf{x}) - \lambda \mathbf{f}(\mathbf{x}) \cdot \nabla V(\mathbf{x}) &\geq 0 \quad \forall \mathbf{x} \in \mathbb{R}^n. \end{aligned}$$

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Nonlinearity

Nonlinearity **37** (2024) 095022 (17pp)

<https://doi.org/10.1088/1361-6544/ad68bb>

The Lorenz system as a gradient-like system

Jeremy P Parker[©]

$$f(x) \cdot \nabla V(x) \geq \|f(x)\|^2$$

arXiv > math > arXiv:2510.14870

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Mathematics > Dynamical Systems

[Submitted on 16 Oct 2025 (v1), last revised 21 Jan 2026 (this version, v2)]

Computation of attractor dimension and maximal sums of Lyapunov exponents using polynomial optimization

Jeremy P Parker, David Goluskin

$$\begin{aligned} d_L \leq \inf_{\substack{B \in [0,1) \\ V \in \mathcal{C}^1}} (k + B) \quad \text{s.t.} \quad & V(x, w, y^{k+1}) \geq |w|^2 \\ & f \cdot D_x V + p_k^B \cdot D_w V + (Df^{[k+1]} y^{k+1}) \cdot D_{y^{k+1}} V \leq 0, \end{aligned}$$

arXiv > math > arXiv:2510.13650

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[Submitted on 15 Oct 2025 (v1), last revised 23 Jun 2026 (this version, v2)]

Computation of minimal periods for ordinary differential equations

Jeremy P. Parker, David Goluskin

$$B^a = \inf_{\substack{Q \succ 0 \\ V \in \mathcal{C}^1(\Omega)}} B \quad \text{s.t.} \quad Ba^T Qa - (\mathcal{L}_f a)^T Q (\mathcal{L}_f a) + \mathcal{L}_f V \geq 0 \quad \forall x \in \Omega.$$



Other Lyapunov/auxiliary/storage function methods

IFAC PapersOnLine 52-16 (2019) 484–489

IEEE TRANSACTIONS ON AUTOMATIC CONTROL, VOL. 59, NO. 2, FEBRUARY 2014

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Using SOS and Sublevel Set Volume Minimization for Estimation of Forward Reachable Sets

Morgan Jones * Matthew M. Peet **

$$\begin{aligned} \min_X \{ & D(X, FR(X_0, f, T, Y, \gamma)) \} \\ \text{subject to: } & X = \{x \in \mathbb{R}^n : V(x, T) \leq 1 + \gamma\} \\ & V(x, 0) \leq 1 \text{ for all } x \in X_0. \\ & \frac{\partial V}{\partial t}(x, t) + \nabla_x V(x, t)^T f(x, u, w) \leq w^T w, \\ & \text{for all } x \in X_c, t \in [0, T], u \in Y, w \in \mathbb{R}^{m_w}. \end{aligned}$$

Proceedings of the
47th IEEE Conference on Decision and Control
Cancun, Mexico, Dec. 9-11, 2008

ThB10.1

Nonlinear optimal control synthesis via occupation measures

Didier Henrion, Jean B. Lasserre and Carlo Savorgnan

$$\frac{\partial \varphi(x, t)}{\partial t} + \nabla_x \varphi(x, t) f(t, x, u) + h(t, x, u) \geq 0 \quad (13)$$

for all $(t, x, u) \in \mathcal{C}_T$ and

$$H(x) - \varphi(T, x) \geq 0 \quad \forall x \in \mathcal{C}_F. \quad (14)$$

Convex Computation of the Region of Attraction of Polynomial Control Systems

Didier Henrion and Milan Korda

$$\begin{aligned} d^* = \inf \int_X & w(x) d\lambda(x) \\ \text{s.t. } & \mathcal{L}v(t, x, u) \leq 0, \quad \forall (t, x, u) \in [0, T] \times X \times U \\ & w(x) \geq v(0, x) + 1, \quad \forall x \in X \\ & v(T, x) \geq 0, \quad \forall x \in X_T \\ & w(x) \geq 0, \quad \forall x \in X \end{aligned} \quad (16)$$



Periodic solutions to ODEs

Consider an autonomous ODE system

$$\frac{d}{dt}x(t) = f(x(t)), \quad f \in C^1(\Omega, \mathbb{R}^n), \quad \Omega \subset \mathbb{R}^n$$

A solution $x(t)$ is a periodic orbit/limit cycle if for some $T > 0$ we have $x(t + T) = x(t)$ for all t but $x(t)$ is **not constant**.

Then we call T a period of the system.



Lower bounds for periods: intuition

Theorem (Fenchel 1929)

The total absolute curvature of a smooth closed curve is at least 2π .

- The local curvature and speed of trajectories is limited, so the periods should be bounded below, at least locally.



Lower bounds for periods: existing results

Theorem (Yorke 1969)

If f is Lipschitz continuous with constant L , i.e. $|f(y) - f(x)| \leq L|y - x|$ for all $x, y \in \Omega$, then periods satisfy

$$T \geq \frac{2\pi}{L}.$$

- This holds for all Hilbert spaces. For Banach spaces $T \geq 6/L$.
- More results for different classes of (infinite dimensional) system.
- This is a sharp bound, e.g. linear oscillator.
- This is not sharp in general for a particular system.
- In fact, f is not generally Lipschitz, but we typically still expect periods to be bounded below.

A lemma

Theorem (Wirtinger's inequality)

If $f \in C^1([0, T], \mathbb{R}^m)$ satisfies $f(T) = f(0)$ and $\int_0^T f(t) dt = 0$, then

$$\int_0^T |f(t)|^2 dt \leq \left(\frac{T}{2\pi}\right)^2 \int_0^T |f'(t)|^2 dt.$$

- Poincaré inequality on the circle
- Equality if the components of f are sinusoidal with period T .

Sketch proof

Theorem (Yorke 1969)

If f is Lipschitz continuous with constant L then periods satisfy

$$T \geq \frac{2\pi}{L}.$$

Proof.

Consider any periodic orbit with period T .

$f(x(t))$ satisfies conditions for Wirtinger's inequality, so

$$\left(\frac{2\pi}{T}\right)^2 \int_0^T |f(x(t))|^2 dt \leq \int_0^T |\mathcal{L}_f f(x(t))|^2 dt.$$

Also $|\mathcal{L}_f f(x)| \leq L|f(x)|$, so

$$\int_0^T |\mathcal{L}_f f(x(t))|^2 dt \leq L^2 \int_0^T |f(x(t))|^2 dt.$$

Generalising

Proposition 1

Suppose there exist $B > 0$, $V \in C^1(\Omega, \mathbb{R})$ and $\Phi \in C^1(\Omega, \mathbb{R}^m)$ such that

$$B |\Phi(x)|^2 - |\mathcal{L}_f \Phi(x)|^2 + \mathcal{L}_f V(x) \geq 0, \quad \forall x \in \Omega.$$

Suppose $x(t)$ is T -periodic for some $T > 0$. If $\Phi(x(t))$ satisfies the mean-zero condition $\int_0^T \Phi(x(t)) dt = 0$ but does not vanish everywhere on $x(t)$, then

$$T \geq \frac{2\pi}{\sqrt{B}}.$$

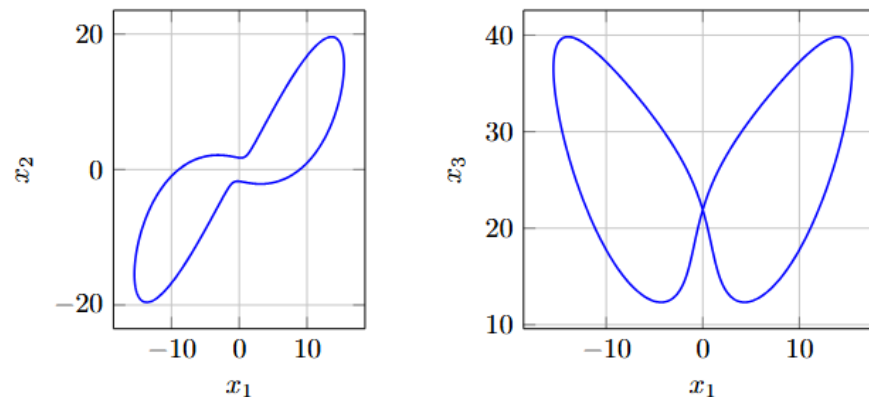
- This reduces to the Yorke result when $\Phi = f$, $V = 0$ and $B = L^2$.



Example: Lorenz '63

$$\begin{aligned}\frac{dx_1}{dt} &= \sigma(x_2 - x_1), \\ \frac{dx_2}{dt} &= x_1(\rho - x_3) - x_2, \\ \frac{dx_3}{dt} &= x_1x_2 - \beta x_3.\end{aligned}$$

$$\beta = \frac{8}{3}, \sigma = 10, \rho = 28$$



Shortest known orbit with $T \approx 1.5587$

Many orbits are invariant under $x \mapsto (-x_1, -x_2, x_3)$.

In that case, x_1 must integrate to zero, and $x_1 \equiv 0$ iff the trajectory tends to the origin.

Example: Lorenz '63

Proposition 1

Suppose there exist $B > 0$, $V \in C^1(\Omega, \mathbb{R})$ and $\Phi \in C^1(\Omega, \mathbb{R}^m)$ such that

$$B |\Phi(x)|^2 - |\mathcal{L}_f \Phi(x)|^2 + \mathcal{L}_f V(x) \geq 0, \quad \forall x \in \Omega.$$

Suppose $x(t)$ is T -periodic for some $T > 0$. If $\Phi(x(t))$ satisfies the mean-zero condition $\int_0^T \Phi(x(t)) dt = 0$ but does not vanish everywhere on $x(t)$, then

$$T \geq \frac{2\pi}{\sqrt{B}}.$$

Take $\Phi(x) = x_1$, $B = \sigma^2(\rho - 1)$, $V = \frac{1}{2}(\sigma(\rho - 2)x_1^2 - \sigma^2 x_2^2 - \sigma^2 x_3^2)$, then

$$B |\Phi(x)|^2 - |\mathcal{L}_f \Phi(x)|^2 + \mathcal{L}_f V(x) = \beta \sigma^2 x_3^2$$

so symmetric orbits have

$$T \geq \frac{2\pi}{\sigma\sqrt{\rho-1}} \approx 0.1209 \ll 1.5587$$

Relaxing to an SDP

Proposition 1

Suppose there exist $B > 0$, $V \in C^1(\Omega, \mathbb{R})$ and $\Phi \in C^1(\Omega, \mathbb{R}^m)$ such that

$$B |\Phi(x)|^2 - |\mathcal{L}_f \Phi(x)|^2 + \mathcal{L}_f V(x) \geq 0, \quad \forall x \in \Omega.$$

Suppose $x(t)$ is T -periodic for some $T > 0$. If $\Phi(x(t))$ satisfies the mean-zero condition $\int_0^T \Phi(x(t)) dt = 0$ but does not vanish everywhere on $x(t)$, then

$$\text{Key idea: choose } \Phi = Ua \text{ for some } a, Q = U^T U \quad T \geq \frac{2\pi}{\sqrt{B}}.$$

Proposition 2

Suppose there exist $B > 0$, $V \in C^1(\Omega, \mathbb{R})$, $a \in C^1(\Omega, \mathbb{R}^m)$ and $Q \succ 0$ such that

$$B a^T Q a - (\mathcal{L}_f a)^T Q (\mathcal{L}_f a) + \mathcal{L}_f V(x) \geq 0, \quad \forall x \in \Omega.$$

Suppose $x(t)$ is T -periodic for some $T > 0$. If $a(x(t))$ satisfies the mean-zero condition $\int_0^T a(x(t)) dt = 0$ but does not vanish everywhere on $x(t)$, then

$$T \geq \frac{2\pi}{\sqrt{B}}.$$



Relaxing to an SDP

Proposition 2

Suppose there exist $B > 0$, $V \in C^1(\Omega, \mathbb{R})$, $a \in C^1(\Omega, \mathbb{R}^m)$ and $Q \succ 0$ such that

$$B a^T Q a - (\mathcal{L}_f a)^T Q (\mathcal{L}_f a) + \mathcal{L}_f V(x) \geq 0, \quad \forall x \in \Omega.$$

Suppose $x(t)$ is T -periodic for some $T > 0$. If $a(x(t))$ satisfies the mean-zero condition $\int_0^T a(x(t)) dt = 0$ but does not vanish everywhere on $x(t)$, then

$$T \geq \frac{2\pi}{\sqrt{B}}.$$

Proposition 3

Suppose $f \in \mathbb{R}^n[x]$ and there exist $B > 0$, $V \in \mathbb{R}[x]$, $w \in \mathbb{R}^m[x]$ and $Q \succ 0$ such that

$$B a^T Q a - (\mathcal{L}_f a)^T Q (\mathcal{L}_f a) + \mathcal{L}_f V(x) \geq 0, \quad \forall x \in \Omega,$$

where $a = L_f w$, and suppose also that the n components of x are among the m components of w .

Then all periods are bounded by

$$T \geq \frac{2\pi}{\sqrt{B}}.$$



Concretely

Suppose

$$\Omega = \{x \in \mathbb{R}^n \mid g_i(x) \geq 0 \text{ for } i = 1, \dots, I, \ h_j(x) = 0 \text{ for } j = 1, \dots, J\},$$

Choose bases $a = \mathcal{L}_f w, b_i, c, c_j$. Let $\Phi = Ua$ and $Q = U^T U$. All periods T are bounded below via

$$\left(\frac{2\pi}{T}\right)^2 \leq \inf_{\substack{Q \in \mathbb{S}^m \\ P_i \in \mathbb{S}^{m_i} \\ v \in \mathbb{R}^{n_0} \\ r_j \in \mathbb{R}^{n_j}}} B \quad \text{s.t.} \quad Ba^T Qa - (\mathcal{L}_f^2 a)^T Q (\mathcal{L}_f^2 a) + v^T \mathcal{L}_f c = b_0^T P_0 b_0 + \sum_{i=1}^I b_i^T P_i b_i g_i + \sum_{j=1}^J r_j^T c_j h_j. \\ Q \succ 0 \text{ and all } P_i \succeq 0.$$

Validation

$$Ba^T Qa - (\mathcal{L}_f^2 a)^T Q (\mathcal{L}_f^2 a) + v^T \mathcal{L}_f c = b_0^T P_0 b_0 + \sum_{i=1}^I b_i^T P_i b_i g_i + \sum_{j=1}^J r_j^T c_j h_j. \quad (20)$$

Algorithm 1 Verified computation of lower bounds on periods using rational arithmetic.

Inputs: $f \in \mathbb{R}^n[x]$, $a \in \mathbb{R}^m[x]$, $c \in \mathbb{R}^{n_0}[x]$, $b_i \in \mathbb{R}^{m_i}[x]$ for $i = 0, \dots, I$, and $c_j \in \mathbb{R}^{n_j}[x]$ for $j = 1, \dots, J$, all with integer coefficients.

Output: $B \in \mathbb{Q}$ such that $\inf_{\mathcal{P}} T \geq \frac{2\pi}{\sqrt{B}}$.

1. Fix a first value of $B \in \mathbb{Q}$.
2. Attempt to solve the SDP feasibility problem (20), using floating point arithmetic and with the constraint $\text{tr } Q = 1$ included to fix the variables' magnitudes.
3. If the SDP solution fails to converge, or if it returns a solution that does not satisfy the constraints within a specified tolerance, increase B and repeat step 2. If the SDP solution returns floating point variables $Q \in \mathbb{S}^m$, $P_i \in \mathbb{S}^{m_i}$, $v \in \mathbb{R}^{n_0}$ and $r_j \in \mathbb{R}^{n_j}$ that nearly satisfy the constraints, including all matrices being positive definite, proceed with rational validation.
4. Express the equality constraint in (20) as a linear relation $Ay = 0$, where the vector y contains all variables in Q , P_i , v and r_j . Choose A to have integer entries, which may require multiplying all terms by the denominator of B .
5. Use rational Gaussian elimination to construct a matrix K whose columns are an exact basis for the kernel of A .
6. Using the floating point result from step 2, construct a rational vector y_0 that approximately solves $Ay = 0$.
7. Find a floating point vector x_0 minimizing $|Kx_0 - y_0|$ using the LSQR algorithm.
8. Approximate x_0 as a rational vector x , so that $y = Kx$ exactly solves $Ay = 0$.
9. Build rational Q , P_i , v and r_j from y that satisfy the equality constraint in (20) by construction.
10. Validate that the rational matrices satisfy $Q \succ 0$ and $P_i \succeq 0$. This is done via LDL^T decomposition, which is exact with rational arithmetic.
11. If validation fails at any step, increase B and return to step 2. If validation succeeds, decrease B and return to step 2. Carry out a line search for the smallest B value that can be validated, terminating when reaching the desired precision in B .

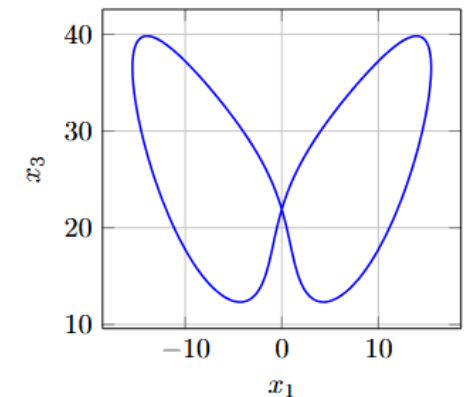
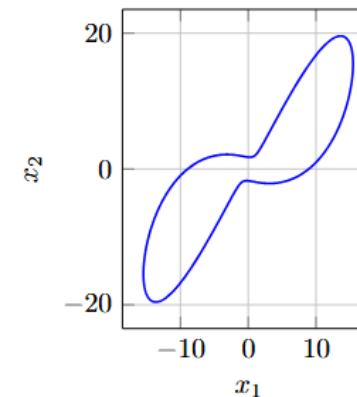
Example: Lorenz '63

Choose w , b and c to be monomials* up to degree d_a , d_b and d_c , and $a = \mathcal{L}_f w$.

$$\Omega = \mathbb{R}^3$$

$$Ba^T Q a - (\mathcal{L}_f^2 a)^T Q (\mathcal{L}_f^2 a) + v^T \mathcal{L}_f c = b_0^T P_0 b_0$$

d_a	d_b	d_c	B	Lower bound on period	SDP time (s)	Validation time (s)
3	3	6	1128	0.5612	0.06	0.03
4	4	8	888	0.6325	0.44	0.62
5	5	10	488	0.8532	4.01	8.98
6	6	12	325	<u>1.0455</u>	29.2	79.7
7	7	14	155.9	<u>1.5096</u>	81.5	573
8	8	16	146.26	<u>1.5586</u>	641	3558



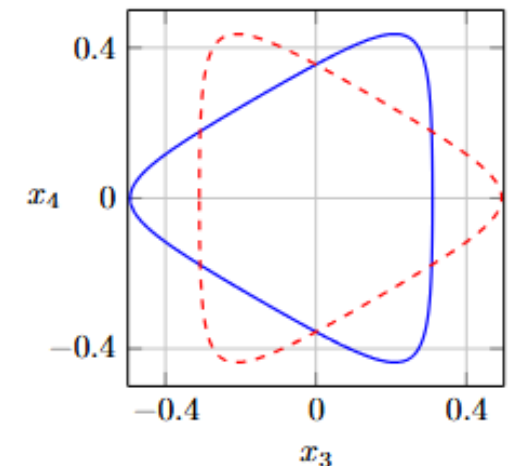
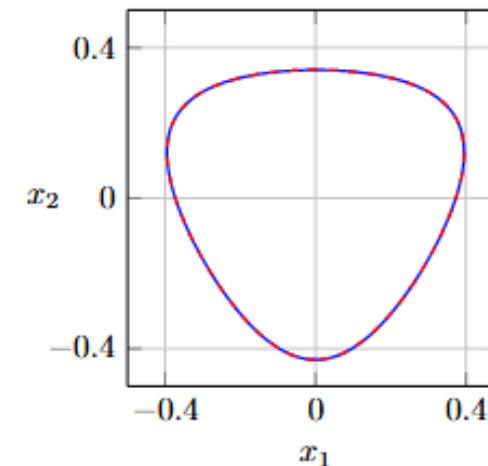
Shortest known orbit with $T \approx 1.55865$

Example: Hénon-Heiles

$$H(x) = \frac{1}{2} (x_1^2 + x_2^2 + x_3^2 + x_4^2) + x_1^2 x_2 - \frac{1}{3} x_2^3, \quad f(x) = (x_3, x_4, -x_1 - 2x_1 x_2, -x_2 - x_1^2 + x_2^2).$$

$$\Omega = \{x \in \mathbb{R}^3 \mid 1 - 6H(x) \geq 0\},$$

d_a	d_b	d_c	B	Lower bound on period	SDP time (s)	Validation time (s)
2	3	5	1.28572	5.5412	0.33	0.01
2	4	7	1.25210	5.6151	7.95	0.38
3	4	7	1.08846	6.0224	8.59	0.49



Shortest known orbits with $T \approx 6.02248$

Summary

- SOS optimisation gives computational global results for polynomial ODE systems
- Can be validated with rational arithmetic to give a computer assisted proof
- We extended SOS for bounds on periods



Things I didn't discuss

- Terrible scaling
- Non-polynomial systems
- Analogous method for polynomial maps
- Extending this to PDEs
- Interval arithmetic validation

